

CAUSAL INFERENCE IN DIRECTED, POSSIBLY CYCLIC, GRAPHICAL MODELS

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Joint work with Elina Robeva ¹

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When Causal Inference meets Statistical Analysis
April 21, 2023

1 THE PROBLEM

2 DISCOVERING THE MARKOV EQUIVALENCE CLASS

- Defining a score
- Greedy optimization of the score

3 DISCOVERING A MARKOV EQUIVALENT GRAPH

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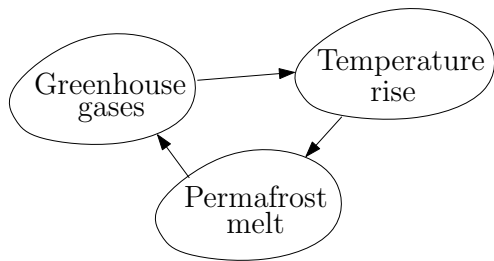
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- Our approach is an extension of the hybrid approach in the acyclic setting [Teyssie and Koller 2005 / Raskutti and Uhler 2018 / ...] to the cyclic setting.
- Previously other algorithms have been proposed which allow for cyclic graphs and have no parametric assumptions [Richardson 2013 / Hyttinen, Hoyer, Eberhardt and Jarvisalo 2013], however, they only output characteristics that are shared by all of the members of the Markov equivalence class of G^* .

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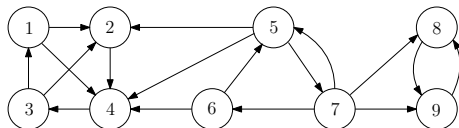


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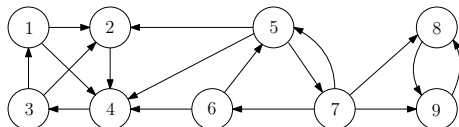


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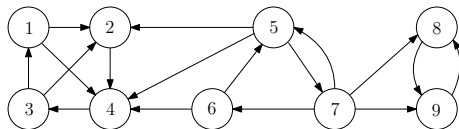


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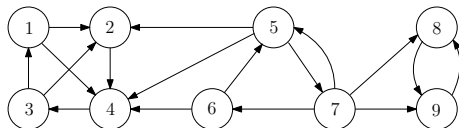


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- the partial order
 - $\{5, 6, 7\} \leq_{G^*} \{1, 2, 3, 4\}$;
 - $\{5, 6, 7\} \leq_{G^*} \{8, 9\}$.

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- Consider the following relation on $\mathcal{SC}(G^*)$: For all $C_1, C_2 \in \mathcal{SC}(G^*)$,

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- $(\mathcal{SC}(G^*), \leq_{G^*})$ is called *the partially ordered partition associated with G^** .
- We find the Markov equivalence class of G^* by optimizing a certain score over the set

$$\mathcal{S} := \left\{ (\mathcal{P}, \pi) \mid \mathcal{P} \subseteq 2^{[n]} \text{ is a partition of } [n] \text{ and } \pi \subseteq \mathcal{P} \times \mathcal{P} \text{ is a partial order on } \mathcal{P}. \right\}.$$

CHARACTERIZATION OF MARKOV EQUIVALENCE

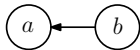
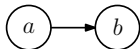
THEOREM [VERMA AND PEARL 1990]

Two DAGs G_1 and G_2 are Markov equivalent if and only if they have the same adjacencies and the same set of immoralities.

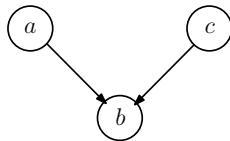
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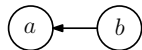
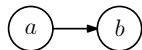


(B) Immorality

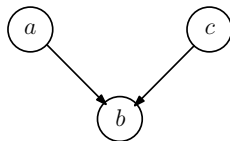
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This characterization does not hold for all directed graphs:

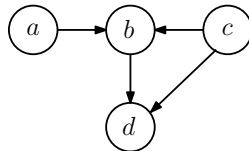
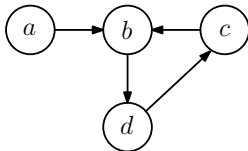


FIGURE: These two graphs have the same adjacencies and immoralities, but $a \not\perp c$ in the first graph, and $a \perp c$ in the second one.

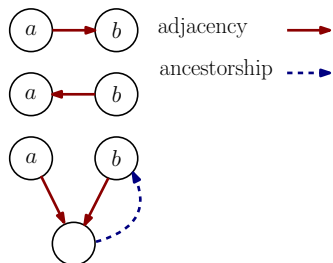
THEOREM [RICHARDSON 1997]

Assume $G_1 = (V, E_1)$ and $G_2 = (V, E_2)$ are two directed graphs. Then G_1 and G_2 are Markov equivalent if and only if the following conditions hold:

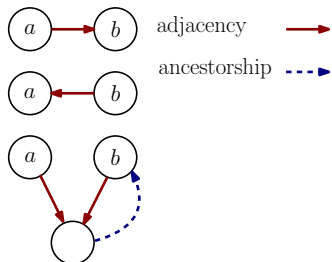
- 1 G_1 and G_2 have the same p-adjacencies.
- 2 G_1 and G_2 have the same set of unshielded non-conductors.
- 3 G_1 and G_2 have the same set of unshielded imperfect non-conductors.
- 4 If (a, b_1, c) and (a, b_2, c) are unshielded imperfect non-conductors (in G_1 and G_2), then b_1 is an ancestor of b_2 in G_1 if and only if b_1 is an ancestor of b_2 in G_2 .
- 5 For any $t \in \mathbb{N}$, triples (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to the uncovered itinerary $P = (a_0, a_1, a_2, \dots, a_{t-1}, a_t, a_{t+1})$ in G_1 if and only if (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to P in G_2 .
- 6 If (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to some uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$ and (a_0, b_1, a_{t+1}) is an unshielded imperfect non-conductor (in G_1 and G_2), then a_1 is an ancestor of b in G_1 if and only if a_1 is an ancestor of b in G_2 .

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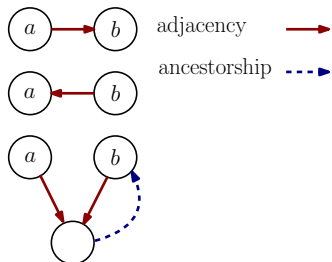


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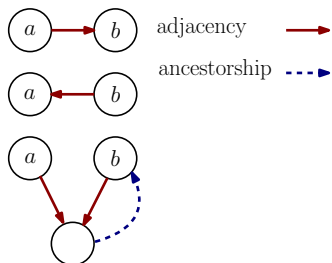
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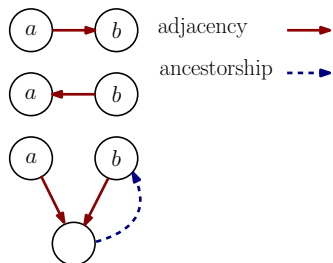
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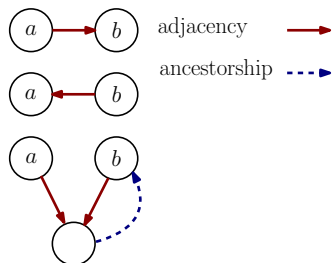
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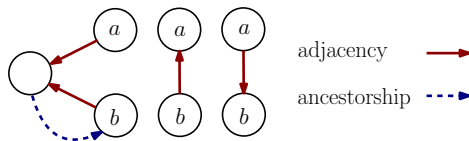


FIGURE: P-adjacency

PROPOSITION

For each $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$E_{(\mathcal{P}, \pi)}^{(1)} := \left\{ (a, b) \in [n]^2 \mid a \neq b, a \not\ll b \mid \bigcup \{ C \in \mathcal{P} \mid C \leq \pi \max\{C_a, \mathcal{P}, C_b, \mathcal{P}\} \} \setminus \{a, b\} \right\}.$$

Also, define $S_1 := \operatorname{argmin}_{(\mathcal{P}, \pi) \in \mathcal{S}} |E_{(\mathcal{P}, \pi)}^{(1)}|$. Then

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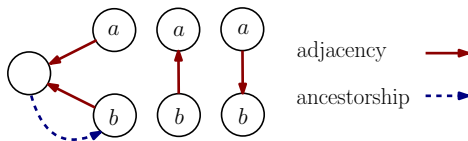


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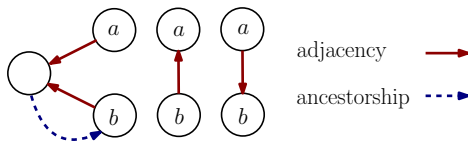


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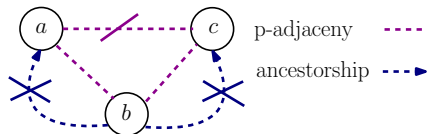
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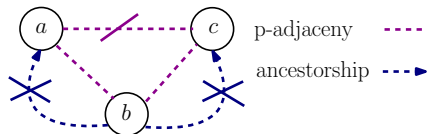
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- for every $(\mathcal{P}, \pi) \in S_1$, the set $E_{(\mathcal{P}, \pi)}^{(1)}$ is equal to the set of p-adjacencies in G^* .

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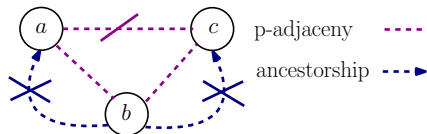


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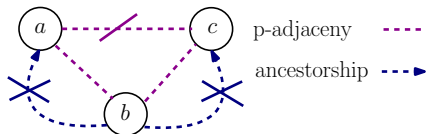
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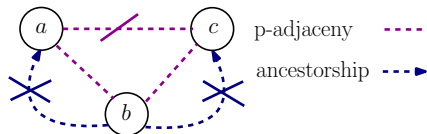
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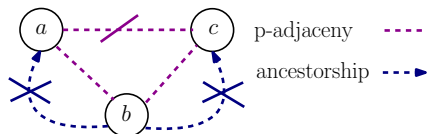
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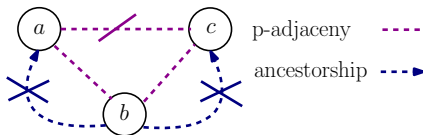


FIGURE: Unshielded non-conductor

PROPOSITION

For each $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$E_{(\mathcal{P}, \pi)}^{(2)} := \{(a, b, c) \in [n]^3 \mid a, b, c \text{ are distinct, } (a, b), (c, b) \in E_{(\mathcal{P}, \pi)}^{(1)}, (a, c) \notin E_{(\mathcal{P}, \pi)}^{(1)}, \\ C_{b, \mathcal{P}} \not\prec_{\pi} C_{a, \mathcal{P}}, C_{b, \mathcal{P}} \not\prec_{\pi} C_{c, \mathcal{P}}\}.$$

Also, define $S_2 := \operatorname{argmax}_{(\mathcal{P}, \pi) \in S_1} |E_{(\mathcal{P}, \pi)}^{(2)}|$. Then

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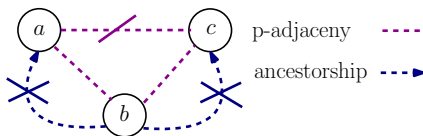


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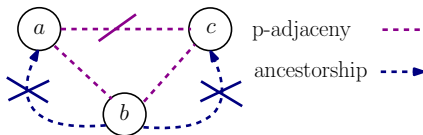


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Also, define $S_2 := \operatorname{argmax}_{(\mathcal{P}, \pi) \in S_1} |E_{(\mathcal{P}, \pi)}^{(2)}|$. Then

- the partially ordered partition associated with G^* is in S_2 , and
- for every $(\mathcal{P}, \pi) \in S_2$, the set $E_{(\mathcal{P}, \pi)}^{(2)}$ is equal to the set of unshielded non-conductors in G^* .

$$E_{(\mathcal{P}, \pi)}^{(3)} := \{(a, b, c) \in [n]^3 \mid (a, b, c) \in E_{(\mathcal{P}, \pi)}^{(2)}, \\ a \perp\!\!\!\perp c \mid \bigcup \{C \in \mathcal{P} \mid C \leq_{\pi} \max\{C_{a, \mathcal{P}}, C_{b, \mathcal{P}}, C_{c, \mathcal{P}}\}\} \setminus \{a, c\}\}.$$

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$$E_{(\mathcal{P}, \pi)}^{(4)} := \{((a, b_1, c), (a, b_2, c)) \in [n]^3 \times [n]^3 \mid (a, b_1, c), (a, b_2, c) \in E_{(\mathcal{P}, \pi)}^{(3)} \\ C_{b_1, \mathcal{P}} \leq_{\pi} C_{b_2, \mathcal{P}}\}.$$

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For all $t \in [n - 2]$,

$$D_{(\mathcal{P}, \pi)}^{(t)} := \{(a_0, a_1, \dots, a_t, a_{t+1}) \in [n]^{t+2} \mid a_0, a_1, \dots, a_{t+1} \text{ are distinct,} \\ (a_i, a_{i-1}) \in E_{(\mathcal{P}, \pi)}^{(1)} \forall i \in [t+1], \\ (a_i, a_j) \notin E_{(\mathcal{P}, \pi)}^{(1)} \forall i \in \{2, \dots, t+1\}, j \leq i-2, \\ C_{a_1, \mathcal{P}} = C_{a_2, \mathcal{P}} = \dots = C_{a_t, \mathcal{P}}, \\ C_{a_1, \mathcal{P}} \not\leq_{\pi} C_{a_0, \mathcal{P}}, C_{a_1, \mathcal{P}} \not\leq_{\pi} C_{a_{t+1}, \mathcal{P}}\}.$$

$$E_{(\mathcal{P}, \pi)}^{(6)} := \bigcup_{t=1}^{n-2} \{ (a_0, a_1, \dots, a_t, a_{t+1}, a_0, b, a_{t+1}) \in [n]^{t+5} \mid (a_0, b, a_{t+1}) \in E_{(\mathcal{P}, \pi)}^{(3)} \}$$

$$(a_0, a_1, \dots, a_t, a_{t+1}) \in D_{(\mathcal{P}, \pi)}^{(t)},$$

$$C_{a_1, \mathcal{P}} \leq_{\pi} C_{b, \mathcal{P}} \}.$$

- For a given set of conditional independence statements and for any $(\mathcal{P}, \pi) \in \mathcal{S}$, we define the *graphical score* of (\mathcal{P}, π) , denoted by $\text{GS}(\mathcal{P}, \pi)$, to be

$$\text{GS}(\mathcal{P}, \pi) := \left(\begin{array}{l} |E_{(\mathcal{P}, \pi)}^{(1)}|, -|E_{(\mathcal{P}, \pi)}^{(2)}|, -|E_{(\mathcal{P}, \pi)}^{(3)}|, |E_{(\mathcal{P}, \pi)}^{(4)}|, \\ -|D_{(\mathcal{P}, \pi)}^{(2)}|, \dots, -|D_{(\mathcal{P}, \pi)}^{(n-2)}|, |E_{(\mathcal{P}, \pi)}^{(6)}| \end{array} \right).$$

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lexicographical order: For $x, y \in \mathbb{Z}^{n+2}$,

$$x < y \iff j = \min \{ i \in [n+2] \mid x_i \neq y_i \} \text{ and } x_j < y_j.$$

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THEOREM

Each minimizer of the graphical score over \mathcal{S} , uniquely determines the Markov equivalence class of G^* .

1 THE PROBLEM

2 DISCOVERING THE MARKOV EQUIVALENCE CLASS

- Defining a score
- Greedy optimization of the score

3 DISCOVERING A MARKOV EQUIVALENT GRAPH

Consider the graph $\mathcal{G} = (\mathcal{S}, \mathcal{E})$, where $((\mathcal{P}_1, \pi_1), (\mathcal{P}_2, \pi_2)) \in \mathcal{E}$ iff $\text{GS}(\mathcal{P}_1, \pi_1) \geq \text{GS}(\mathcal{P}_2, \pi_2)$ and

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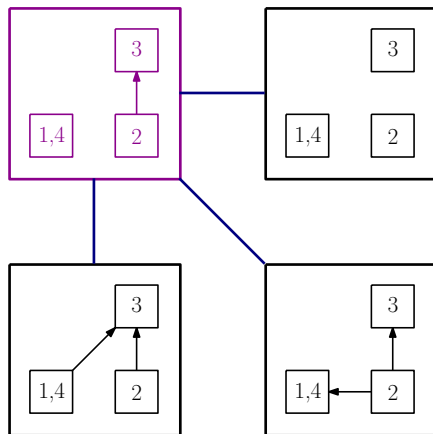
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GRAPH OF THE PARTIALLY ORDERED PARTITIONS

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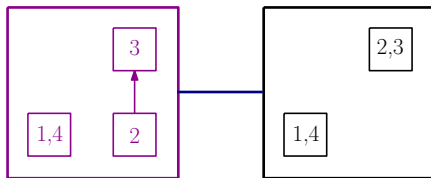
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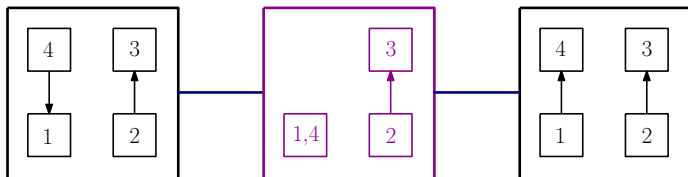
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Based on experimental evidence, we conjecture that for any $(\mathcal{P}, \pi) \in \mathcal{S}$, there exist an optimal partially ordered partition (\mathcal{P}_0, π_0) and a directed path in \mathcal{G} such that the path starts from (\mathcal{P}, π) and ends at (\mathcal{P}_0, π_0) .

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Algorithm 1 Markov equivalence class discovery

Input: The set of all the conditional independence statements satisfied by distribution \mathbb{P} and an initial $(\mathcal{P}_1, \pi_1) \in \mathcal{S}$.

Output: An optimal partially ordered partition.

- 1: Set $\hat{\mathcal{P}} := \mathcal{P}_1$ and $\hat{\pi} := \pi_1$.
 - 2: Perform a depth-first search on \mathcal{G} with root $(\hat{\mathcal{P}}, \hat{\pi})$ to find a directed path from $(\hat{\mathcal{P}}, \hat{\pi})$ to a partially ordered partition $(\tilde{\mathcal{P}}, \tilde{\pi})$ with $\text{GS}(\hat{\mathcal{P}}, \hat{\pi}) > \text{GS}(\tilde{\mathcal{P}}, \tilde{\pi})$.
 - 3: **if** $(\tilde{\mathcal{P}}, \tilde{\pi})$ is found **then**
 - 4: Set $\hat{\mathcal{P}} := \tilde{\mathcal{P}}$ and $\hat{\pi} := \tilde{\pi}$, and go back to step 2.
 - 5: **else**
 - 6: **return** $(\hat{\mathcal{P}}, \hat{\pi})$.
 - 7: **end if**
-

- Once this algorithm finds an optimal partially ordered partition $(\hat{\mathcal{P}}, \hat{\pi})$, it needs to perform a full depth-first search on \mathcal{G} with root $(\hat{\mathcal{P}}, \hat{\pi})$ before it makes sure that there is no partially ordered partition with a lower score.

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- One can choose to repeat the algorithm M times with M different initial partially ordered partitions to get better results, where M is also part of the input.

- Random graphs were generated according to the Erdős–Rényi model and using the R library `igraph`. The set of all d-separations satisfied by each graph was then generated and given to this algorithm as part of its input.

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- We used the following initial partially ordered partitions ($M = 3$):

$$\mathcal{P}_1 = \{\{1, \dots, n\}\},$$

$$\mathcal{P}_2 = \left\{ \left\{ 1, \dots, \lfloor \frac{n}{2} \rfloor \right\}, \left\{ \lfloor \frac{n}{2} \rfloor + 1, \dots, n \right\} \right\}, \left\{ 1, \dots, \lfloor \frac{n}{2} \rfloor \right\} \leq_{\pi_2} \left\{ \lfloor \frac{n}{2} \rfloor + 1, \dots, n \right\},$$

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- The positive integers N used in our experiments were the following:

number of vertices of the graph (n)	N
7	30
8	30
9	40
10	50

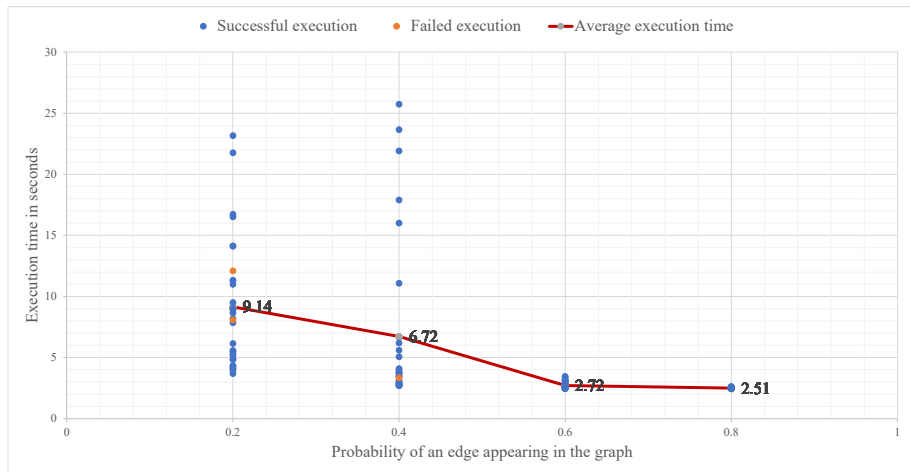


FIGURE: Results of running the greedy optimization algorithm on 120 random graphs with 7 vertices and different levels of sparsity. Thirty graphs were tested for each of the probabilities 0.2, 0.4, 0.6, and 0.8. From left to right, the success rates are 0.93, 0.97, 1 and 1.



FIGURE: Results of running the greedy optimization algorithm on 120 random graphs with different numbers of vertices n where each edge appears in the graph with probability 0.3. Thirty graphs were tested for each n . From left to right, the success rates are 0.90, 0.93, 0.97 and 0.93.

1 THE PROBLEM

2 DISCOVERING THE MARKOV EQUIVALENCE CLASS

- Defining a score
- Greedy optimization of the score

3 DISCOVERING A MARKOV EQUIVALENT GRAPH

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- 3 If $(a, b, c) \in E_{(\mathcal{P}, \pi)}^{(2)} \setminus E_{(\mathcal{P}, \pi)}^{(3)}$ and for all $b' \in [n]$ with $C_{b', \mathcal{P}} <_{\pi} C_{b, \mathcal{P}}$ and $(a, b', c) \in E_{(\mathcal{P}, \pi)}^{(2)}$, $(a, b', c) \in E_{(\mathcal{P}, \pi)}^{(3)}$, then a and c have a common child in $C_{b, \mathcal{P}}$ in graph G .

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Then G is Markov equivalent to G^* .

Moreover, G^* satisfies properties 1 to 4 with respect to the partially ordered partition associated with G^* .

Algorithm 3 Markov equivalent graph discovery

Input: The set of all the conditional independence statements satisfied by distribution \mathbb{P} , an optimal partially ordered partition (\mathcal{P}, π) .

Output: A graph Markov equivalent to G^* .

- 1: Set $\hat{\mathcal{P}} := \mathcal{P}$ and $\hat{\pi} := \pi$.
- 2: **if** there exist $a, b \in [n]$ such that $(a, b) \in E_{(\hat{\mathcal{P}}, \hat{\pi})}^{(1)}$, but $C_{a, \hat{\mathcal{P}}} \not\prec_{\hat{\pi}} C_{b, \hat{\mathcal{P}}}$ and $C_{b, \hat{\mathcal{P}}} \not\prec_{\hat{\pi}} C_{a, \hat{\mathcal{P}}}$ **then**
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SCCR ALGORITHM SIMULATIONS

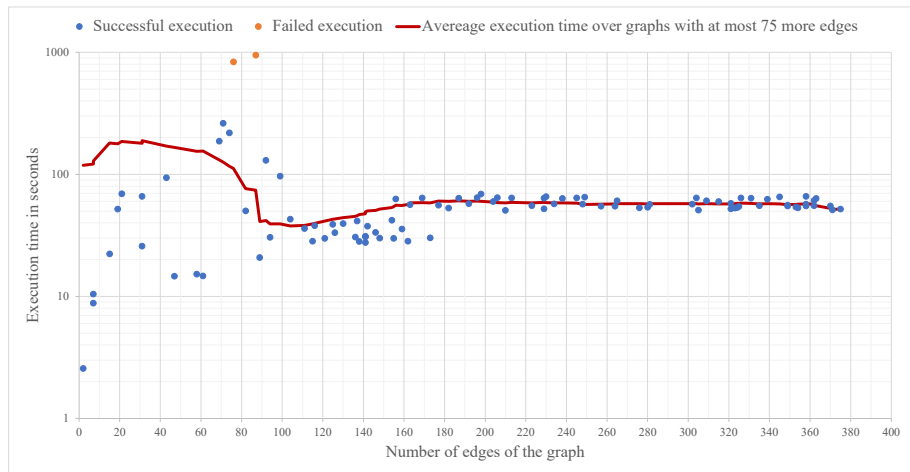


FIGURE: Results of running the SCCR algorithm with $N = 100$ on the partially ordered partitions associated with 100 random graphs with 20 vertices generated according to the Erdős–Rényi model. The success rate is 0.98 and the average execution time is 71.90 seconds.

- Random graphs were generated according to the Erdős–Rényi model and using the R library `igraph`. The set of all d-separations satisfied by each graph was then generated and given to this algorithm as its input.

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- We used the SCCR algorithm with $N = 100$.
- In the Markov equivalent graph discovery algorithm, we restricted ourselves to testing at most 300 optimal partially ordered partitions.



FIGURE: Results of running the Markov equivalent graph discovery algorithm on 120 random graphs with 7 vertices and different levels of sparsity. Thirty graphs were tested for each of the probabilities 0.2, 0.4, 0.6, and 0.8. From left to right, the success rates are 0.3, 0.97, 1 and 1.

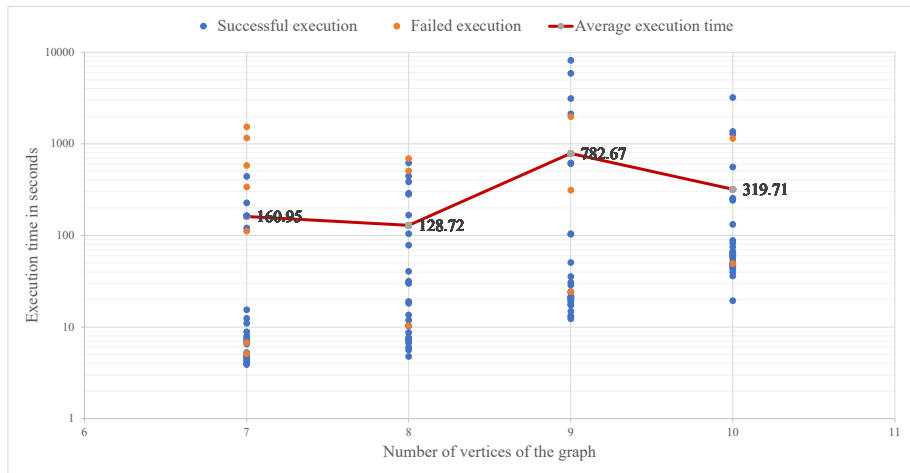


FIGURE: Results of running the Markov equivalent graph discovery algorithm on 120 random graphs with different numbers of vertices n where each edge appears in the graph with probability 0.3. Thirty graphs were tested for each n . From left to right, the success rates are 0.77, 0.90, 0.90 and 0.93.

Thank you for listening!



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- 5 OTHER CONDITIONS OF RICHARDSON'S EQUIVALENCE THEOREM
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- 8 SCCR ALGORITHM

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- **The propensity to consume:** If the group of all consumers in the society are provided with the total income r_t in year t , then they will spend a total amount u_t for consumption in that year, equal to

$$u_t = \alpha \cdot r_t + \beta + \varepsilon_t^{(1)},$$

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DIRECTED CYCLES IN THE CAUSAL GRAPH OF AN SCM

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- **Closed market identity:** In a closed market, all the income will either be spent on consumption goods or invested. So,

$$r_t = u_t + v_t.$$

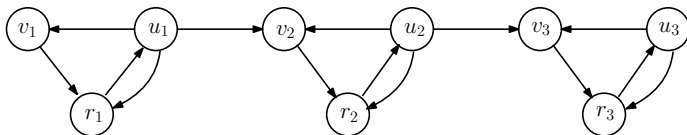
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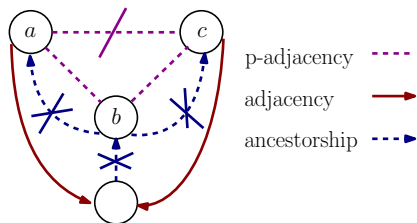
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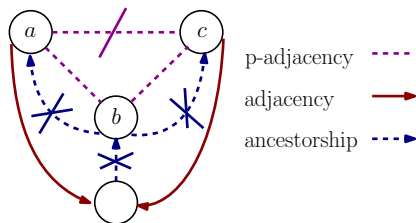
- 4 AN EXAMPLE OF A DIRECTED CYCLE IN THE CAUSAL GRAPH
- 5 OTHER CONDITIONS OF RICHARDSON'S EQUIVALENCE THEOREM**
- 6 GREEDY OPTIMIZATION ALGORITHM
- 7 OPTIMIZATION SIMULATIONS
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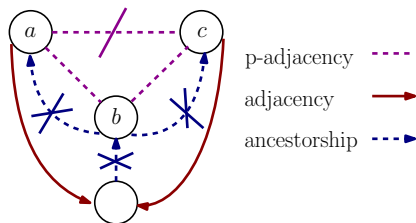


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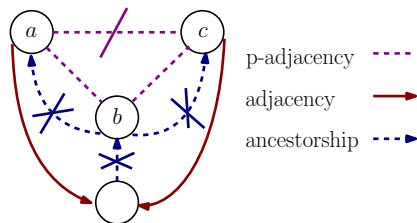
Let $G = (V, E)$ be a directed graph. A triple $(a, b, c) \in V^3$ is said to be an *unshielded imperfect non-conductor* if

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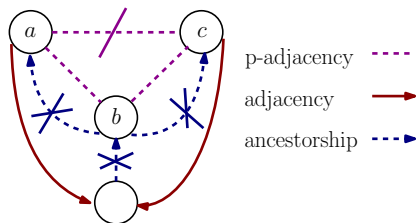
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- b is not the descendant of any of the common children of a and c .

CONDITION 3

G_1 and G_2 have the same set of unshielded imperfect non-conductors.

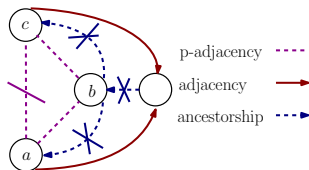


FIGURE: Unshielded imperfect non-conductor

PROPOSITION

For each $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$E_{(\mathcal{P}, \pi)}^{(3)} := \{(a, b, c) \in [n]^3 \mid (a, b, c) \in E_{(\mathcal{P}, \pi)}^{(2)},$$

$$a \perp\!\!\!\perp c \mid \bigcup \{C \in \mathcal{P} \mid C \leq_{\pi} \max\{C_{a, \mathcal{P}}, C_{b, \mathcal{P}}, C_{c, \mathcal{P}}\}\} \setminus \{a, c\}\}.$$

Also, define $S_3 := \operatorname{argmax}_{(\mathcal{P}, \pi) \in \mathcal{S}_2} |E_{(\mathcal{P}, \pi)}^{(3)}|$. Then

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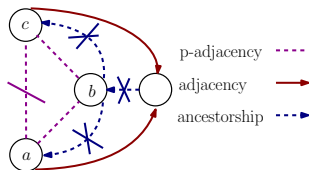


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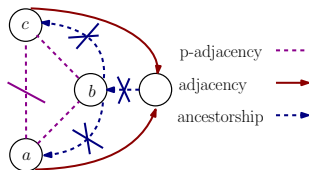


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- the partially ordered partition associated with G^* is in S_3 , and
- for every $(\mathcal{P}, \pi) \in S_3$, the set $E_{(\mathcal{P}, \pi)}^{(3)}$ is equal to the set of unshielded imperfect non-conductors in G^* .

If (a, b_1, c) and (a, b_2, c) are unshielded imperfect non-conductors (in G_1 and G_2), then b_1 is an ancestor of b_2 in G_1 iff b_1 is an ancestor of b_2 in G_2 .

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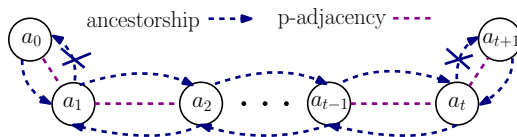
- the partially ordered partition associated with G^* is in S_4 , and
- for every $(\mathcal{P}, \pi) \in S_4$, $((a, b_1, c), (a, b_2, c)) \in E_{(\mathcal{P}, \pi)}^{(4)}$ if and only if $(a, b_1, c), (a, b_2, c)$ are unshielded imperfect non-conductors and b_1 is an ancestor of b_2 in G^* .

CONDITION 5

For any $t \in \mathbb{N}$, triples (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to the uncovered itinerary $P = (a_0, a_1, a_2, \dots, a_{t-1}, a_t, a_{t+1})$ in G_1 iff (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to P in G_2 .

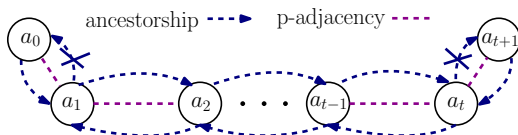
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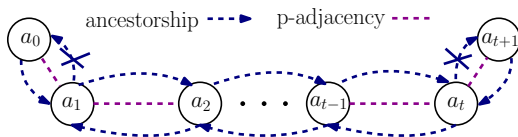
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Let $G = (V, E)$ be a directed graph. For $t \in \mathbb{N}$ and $a_0, a_1, \dots, a_t, a_{t+1} \in V$, the triples $(a_0, a_1, a_2), (a_{t-1}, a_t, a_{t+1})$ are said to be *mutually exclusive with respect to an uncovered itinerary* $(a_0, a_1, a_2, \dots, a_{t-1}, a_t, a_{t+1})$ if

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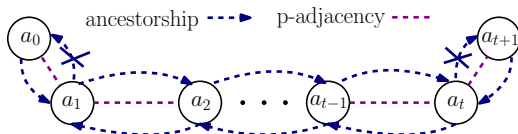


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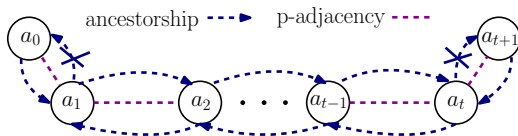


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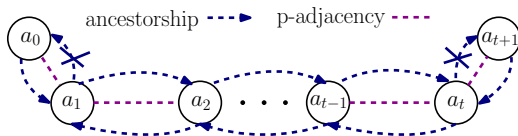


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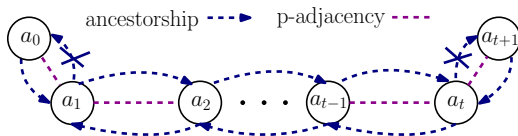


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- a_0 and a_{t+1} are not in this component, and
- a_0 and a_{t+1} are ancestors of a_1 .

PROPOSITION

For all $t \in [n - 2]$ and $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$D_{(\mathcal{P}, \pi)}^{(t)} := \left\{ (a_0, a_1, \dots, a_t, a_{t+1}) \in [n]^{t+2} \mid \begin{array}{l} a_0, a_1, \dots, a_{t+1} \text{ are distinct,} \\ (a_i, a_{i-1}) \in E_{(\mathcal{P}, \pi)}^{(1)} \forall i \in [t+1], \\ (a_i, a_j) \notin E_{(\mathcal{P}, \pi)}^{(1)} \forall i \in \{2, \dots, t+1\}, j \leq i-2, \\ C_{a_1, \mathcal{P}} = C_{a_2, \mathcal{P}} = \dots = C_{a_t, \mathcal{P}}, \\ C_{a_1, \mathcal{P}} \not\leq \pi C_{a_0, \mathcal{P}}, C_{a_1, \mathcal{P}} \not\leq \pi C_{a_{t+1}, \mathcal{P}} \end{array} \right\}.$$

Also, define $S_5^{(1)} := S_4$ and for all $t \in \{2, \dots, n - 2\}$, $S_5^{(t)} := \operatorname{argmax}_{(\mathcal{P}, \pi) \in S_5^{(t-1)}} \left| D_{(\mathcal{P}, \pi)}^{(t)} \right|$.

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PROPOSITION

For all $t \in [n - 2]$ and $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$D_{(\mathcal{P}, \pi)}^{(t)} := \left\{ (a_0, a_1, \dots, a_t, a_{t+1}) \in [n]^{t+2} \mid \begin{array}{l} a_0, a_1, \dots, a_{t+1} \text{ are distinct,} \\ (a_i, a_{i-1}) \in E_{(\mathcal{P}, \pi)}^{(1)} \forall i \in [t+1], \\ (a_i, a_j) \notin E_{(\mathcal{P}, \pi)}^{(1)} \forall i \in \{2, \dots, t+1\}, j \leq i-2, \\ C_{a_1, \mathcal{P}} = C_{a_2, \mathcal{P}} = \dots = C_{a_t, \mathcal{P}}, \\ C_{a_1, \mathcal{P}} \not\leq \pi C_{a_0, \mathcal{P}}, C_{a_1, \mathcal{P}} \not\leq \pi C_{a_{t+1}, \mathcal{P}} \end{array} \right\}.$$

Also, define $S_5^{(1)} := S_4$ and for all $t \in \{2, \dots, n - 2\}$, $S_5^{(t)} := \operatorname{argmax}_{(\mathcal{P}, \pi) \in S_5^{(t-1)}} |D_{(\mathcal{P}, \pi)}^{(t)}|$.

Then

- the partially ordered partition associated with G^* is in $S_5^{(t)}$ for all $t \in [n - 2]$, and
- for all $t \in [n - 2]$ and $(\mathcal{P}, \pi) \in S_5^{(t)}$, $(a_0, a_1, \dots, a_t, a_{t+1}) \in D_{(\mathcal{P}, \pi)}^{(t)}$ if and only if (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to the uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$ in G^* .

CONDITION 6

If (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to some uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$ and (a_0, b_1, a_{t+1}) is an unshielded imperfect non-conductor (in G_1 and G_2), then a_1 is an ancestor of b in G_1 iff a_1 is an ancestor of b in G_2 .

CONDITION 6

If (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to some uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$ and (a_0, b_1, a_{t+1}) is an unshielded imperfect non-conductor (in G_1 and G_2), then a_1 is an ancestor of b in G_1 iff a_1 is an ancestor of b in G_2 .

PROPOSITION

For every $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$E_{(\mathcal{P}, \pi)}^{(6)} := \bigcup_{t=1}^{n-2} \{ ((a_0, a_1, \dots, a_t, a_{t+1}), (a_0, b, a_{t+1})) \in [n]^{t+5} \mid (a_0, b, a_{t+1}) \in E_{(\mathcal{P}, \pi)}^{(3)} \}$$

$$(a_0, a_1, \dots, a_t, a_{t+1}) \in D_{(\mathcal{P}, \pi)}^{(t)},$$

$$C_{a_1, \mathcal{P}} \leq_{\pi} C_{b, \mathcal{P}} \}.$$

Also, define $S_6 := \operatorname{argmin}_{(\mathcal{P}, \pi) \in S_5^{(n-2)}} |E_{(\mathcal{P}, \pi)}^{(6)}|$. Then

CONDITION 6

If (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to some uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$ and (a_0, b_1, a_{t+1}) is an unshielded imperfect non-conductor (in G_1 and G_2), then a_1 is an ancestor of b in G_1 iff a_1 is an ancestor of b in G_2 .

PROPOSITION

For every $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$E_{(\mathcal{P}, \pi)}^{(6)} := \bigcup_{t=1}^{n-2} \{ ((a_0, a_1, \dots, a_t, a_{t+1}), (a_0, b, a_{t+1})) \in [n]^{t+5} \mid (a_0, b, a_{t+1}) \in E_{(\mathcal{P}, \pi)}^{(3)} \}$$

$$(a_0, a_1, \dots, a_t, a_{t+1}) \in D_{(\mathcal{P}, \pi)}^{(t)},$$

$$C_{a_1, \mathcal{P}} \leq \pi C_{b, \mathcal{P}} \}.$$

Also, define $S_6 := \operatorname{argmin}_{(\mathcal{P}, \pi) \in S_5^{(n-2)}} |E_{(\mathcal{P}, \pi)}^{(6)}|$. Then

- the partially ordered partition associated with G^* is in S_6 , and

If (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to some uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$ and (a_0, b_1, a_{t+1}) is an unshielded imperfect non-conductor (in G_1 and G_2), then a_1 is an ancestor of b in G_1 iff a_1 is an ancestor of b in G_2 .

PROPOSITION

For every $(\mathcal{P}, \pi) \in \mathcal{S}$, let

$$E_{(\mathcal{P}, \pi)}^{(6)} := \bigcup_{t=1}^{n-2} \{ ((a_0, a_1, \dots, a_t, a_{t+1}), (a_0, b, a_{t+1})) \in [n]^{t+5} \mid (a_0, b, a_{t+1}) \in E_{(\mathcal{P}, \pi)}^{(3)} \}$$

$$(a_0, a_1, \dots, a_t, a_{t+1}) \in D_{(\mathcal{P}, \pi)}^{(t)},$$

$$C_{a_1, \mathcal{P}} \leq_{\pi} C_{b, \mathcal{P}} \}.$$

Also, define $S_6 := \operatorname{argmin}_{(\mathcal{P}, \pi) \in S_5^{(n-2)}} |E_{(\mathcal{P}, \pi)}^{(6)}|$. Then

- the partially ordered partition associated with G^* is in S_6 , and
- for every $(\mathcal{P}, \pi) \in S_6$, $((a_0, a_1, \dots, a_t, a_{t+1}), (a_0, b, a_{t+1})) \in E_{(\mathcal{P}, \pi)}^{(6)}$ if and only if (a_0, a_1, a_2) and (a_{t-1}, a_t, a_{t+1}) are mutually exclusive with respect to the uncovered itinerary $(a_0, a_1, \dots, a_{t+1})$, (a_0, b_1, a_{t+1}) is an unshielded imperfect non-conductor, and a_1 is an ancestor of b in G^* .

- 4 AN EXAMPLE OF A DIRECTED CYCLE IN THE CAUSAL GRAPH
- 5 OTHER CONDITIONS OF RICHARDSON'S EQUIVALENCE THEOREM
- 6 GREEDY OPTIMIZATION ALGORITHM**
- 7 OPTIMIZATION SIMULATIONS
- 8 SCCR ALGORITHM

Algorithm 2 Greedy Markov equivalence class discovery

Input: The set of all the conditional independence statements satisfied by distribution \mathbb{P} , two positive integers N and M , and initial partially ordered partitions $(\mathcal{P}_1, \pi_1), \dots, (\mathcal{P}_M, \pi_M)$.

Output: An optimal partially ordered partition.

- 1: Set $A := \emptyset$.
 - 2: **for** i in $1 : M$ **do**
 - 3: Set $\hat{\mathcal{P}} := \mathcal{P}_i$ and $\hat{\pi} := \pi_i$.
 - 4: Perform a depth-first search on \mathcal{G} with root $(\hat{\mathcal{P}}, \hat{\pi})$ to find a directed path from $(\hat{\mathcal{P}}, \hat{\pi})$ to a partially ordered partition $(\tilde{\mathcal{P}}, \tilde{\pi})$ with $\text{GS}(\hat{\mathcal{P}}, \hat{\pi}) > \text{GS}(\tilde{\mathcal{P}}, \tilde{\pi})$. Stop the depth-first search once a directed path of length N consisting of partially ordered partitions of score $\text{GS}(\hat{\mathcal{P}}, \hat{\pi})$ is generated.
 - 5: **if** $(\tilde{\mathcal{P}}, \tilde{\pi})$ is found **then**
 - 6: Set $\hat{\mathcal{P}} := \tilde{\mathcal{P}}$ and $\hat{\pi} := \tilde{\pi}$, and go back to step 4.
 - 7: **else**
 - 8: Set $A := A \cup \{(\hat{\mathcal{P}}, \hat{\pi})\}$.
 - 9: **end if**
 - 10: **end for**
 - 11: **return** $\operatorname{argmin}_{(\mathcal{P}, \pi) \in A} \text{GS}(\mathcal{P}, \pi)$.
-

- 4 AN EXAMPLE OF A DIRECTED CYCLE IN THE CAUSAL GRAPH
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FIGURE: Results of running the greedy optimization algorithm on 120 random graphs with different numbers of vertices n where each edge appears in the graph with probability 0.2. Thirty graphs were tested for each n . From left to right, the success rates are 0.93, 0.93, 0.93 and 0.87.

- 4 AN EXAMPLE OF A DIRECTED CYCLE IN THE CAUSAL GRAPH
- 5 OTHER CONDITIONS OF RICHARDSON'S EQUIVALENCE THEOREM
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- 8 SCCR ALGORITHM**

$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

$$\begin{aligned}C &:= \{1, 2, 3, 4, 5, 6\}, \\A_C &:= \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}, \\B_C &:= \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}, \\ComCh_C &:= \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.\end{aligned}$$

Two edges $(a, b), (c, b) \in [n] \times C$ are said to be *incompatible* if $a \neq c$ and one of the following happens:

$$\begin{aligned}C &:= \{1, 2, 3, 4, 5, 6\}, \\A_C &:= \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}, \\B_C &:= \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}, \\ComCh_C &:= \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.\end{aligned}$$

Two edges $(a, b), (c, b) \in [n] \times C$ are said to be *incompatible* if $a \neq c$ and one of the following happens:

- $a, c \in [n] \setminus C$ and $(a, c) \in NoComCh_C$, or

$$\begin{aligned}
 C &:= \{1, 2, 3, 4, 5, 6\}, \\
 A_C &:= \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}, \\
 B_C &:= \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}, \\
 ComCh_C &:= \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.
 \end{aligned}$$

Two edges $(a, b), (c, b) \in [n] \times C$ are said to be *incompatible* if $a \neq c$ and one of the following happens:

- $a, c \in [n] \setminus C$ and $(a, c) \in NoComCh_C$, or
- $C \cap \{a, c\} \neq \emptyset$, $(a, c) \notin A_C \cup B_C$, and $(c, a) \notin A_C \cup B_C$.

$$\begin{aligned}
 C &:= \{1, 2, 3, 4, 5, 6\}, \\
 A_C &:= \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}, \\
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 \end{aligned}$$

Two edges $(a, b), (c, b) \in [n] \times C$ are said to be *incompatible* if $a \neq c$ and one of the following happens:

- $a, c \in [n] \setminus C$ and $(a, c) \in NoComCh_C$, or
- $C \cap \{a, c\} \neq \emptyset$, $(a, c) \notin A_C \cup B_C$, and $(c, a) \notin A_C \cup B_C$.

An edge is said to be *safe* to be added to a set of edges if it's not incompatible with any of them.

AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

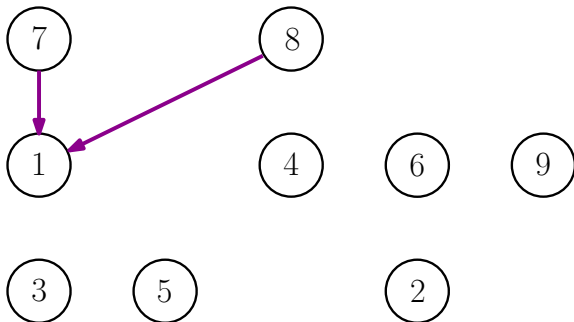
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

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$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

First assigns a common child
to each pair in $ComCh_C$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

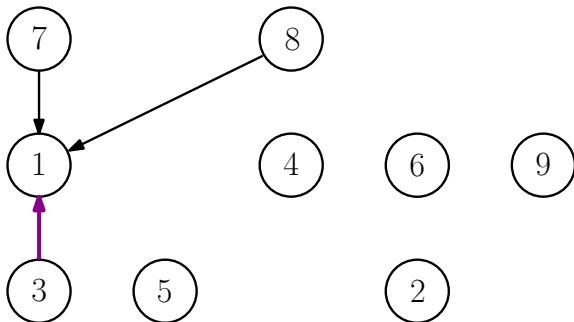
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$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Starts from a vertex in C
with the lowest degree in
 $([n], A_C \cup B_C)$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

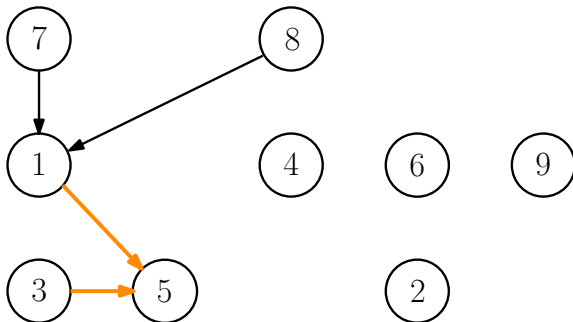
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$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

When an edge is added, the algorithm makes sure it's safe. If not, the algorithm initiates a correction process. In this process, the algorithm first tries to remove the unsafe edge.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

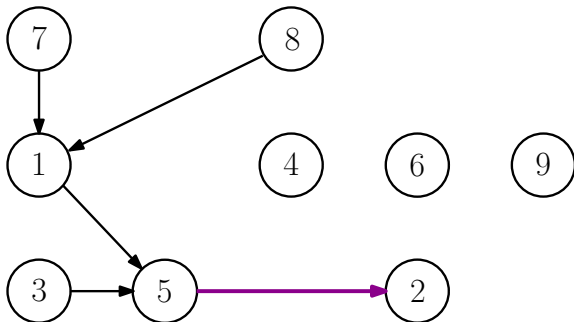
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$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

In each iteration, the algorithm adds an edge from $A_C \cup B_C$ to the construction such that the added edges form an *almost directed path*.



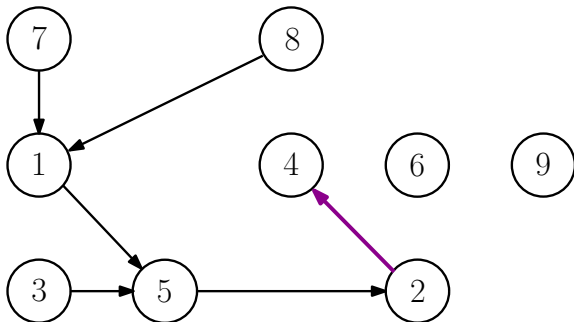
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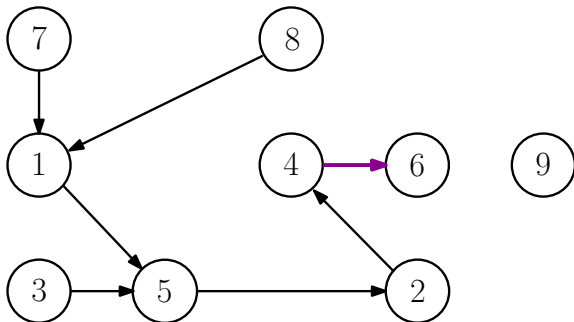
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

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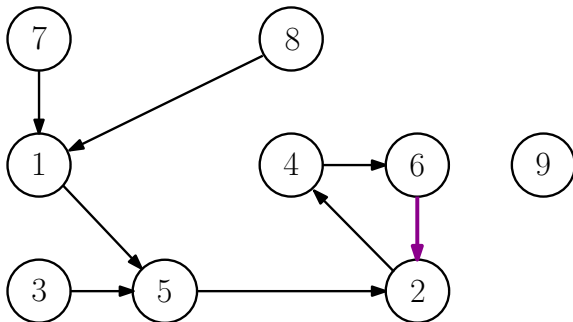
$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

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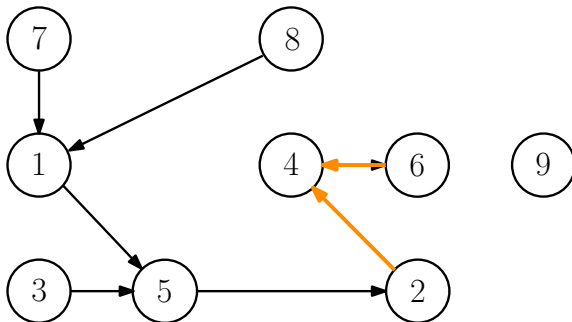
Unsafe edge added
(incompatible with (5, 2))!
Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
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$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Problem resolved!



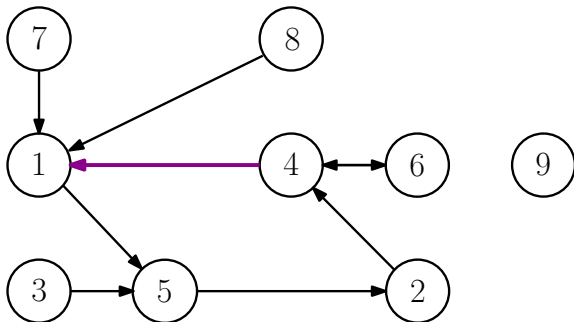
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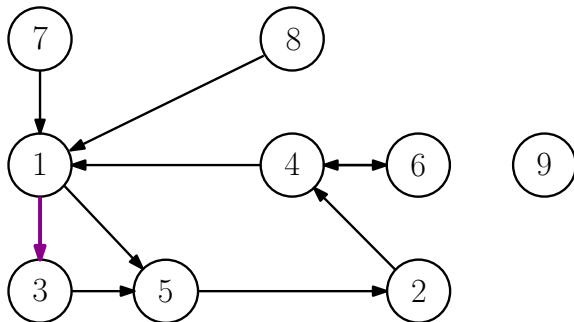
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AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

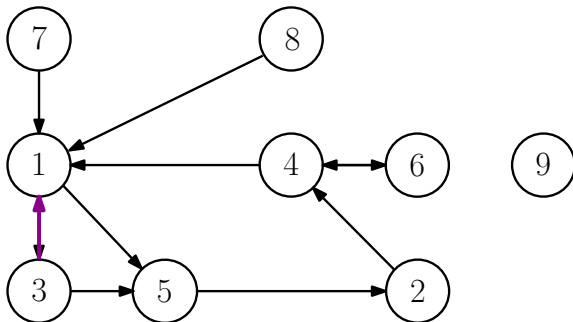
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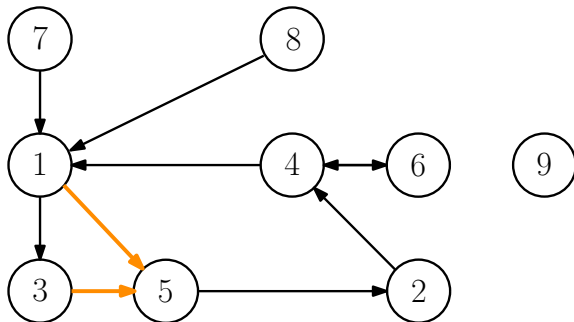
Unsafe edge added
(incompatible with (7, 1))!
Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
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Problem resolved!



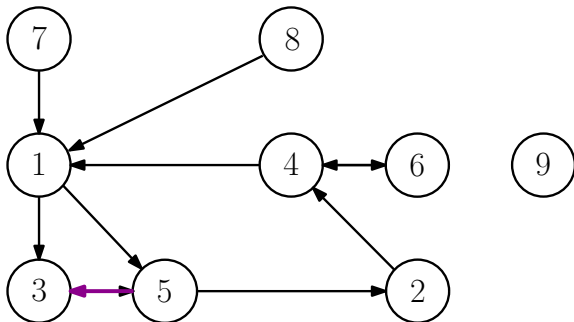
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$C := \{1, 2, 3, 4, 5, 6\}$,

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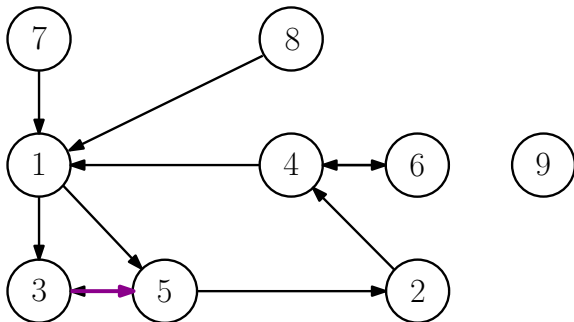
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AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

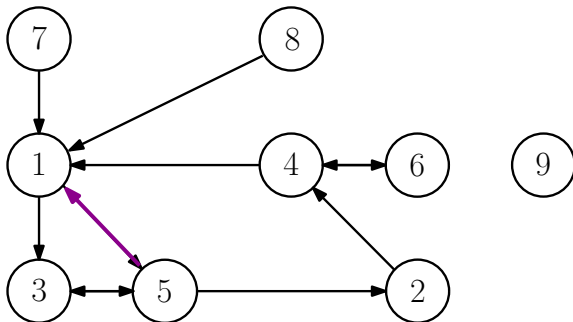
$$C := \{1, 2, 3, 4, 5, 6\},$$

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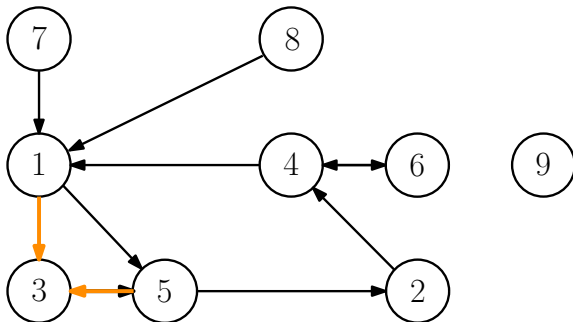
Unsafe edge added
(incompatible with $(7, 1)$)!
Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
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$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Problem resolved!



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

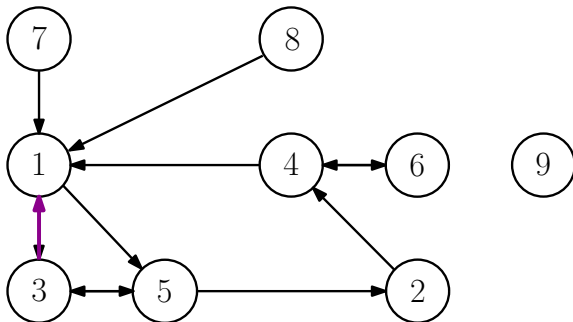
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

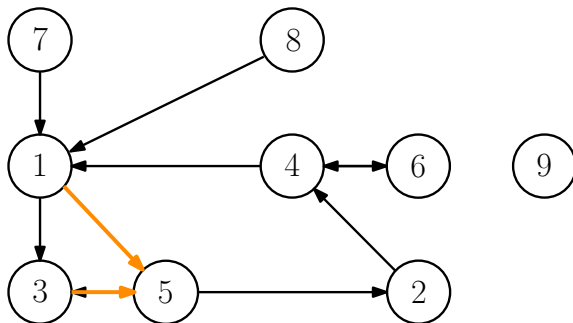
Unsafe edge added
(incompatible with (7, 1))!
Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$
$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Problem resolved!



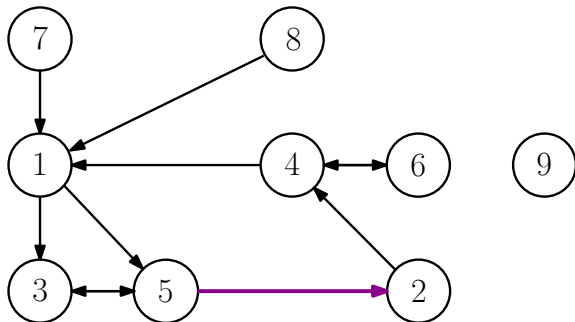
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

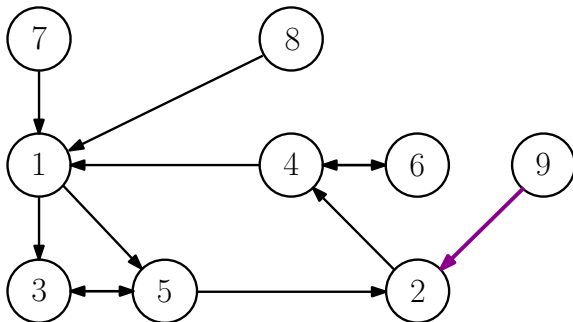
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

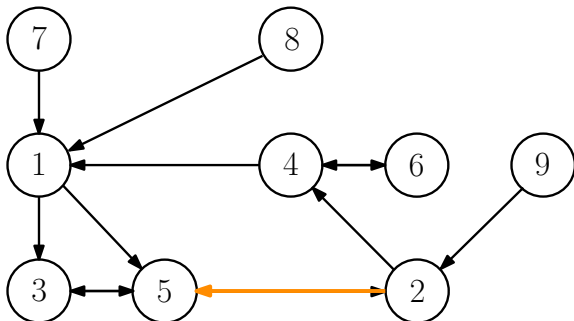
Unsafe edge added
(incompatible with (5, 2))!
Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$\begin{aligned}C &:= \{1, 2, 3, 4, 5, 6\}, \\A_C &:= \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}, \\B_C &:= \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}, \\ComCh_C &:= \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.\end{aligned}$$

In the correction process, when removing fails, the algorithm flips the edge (in case the head and the tail are both in C). Each edge is allowed to be corrected at most once. So, if the flipped edge is incompatible with an edge already affected in a correction process (in this case, $(1, 5)$), the algorithm jumps back to the stage right before the iteration involving the correction of this edge started.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

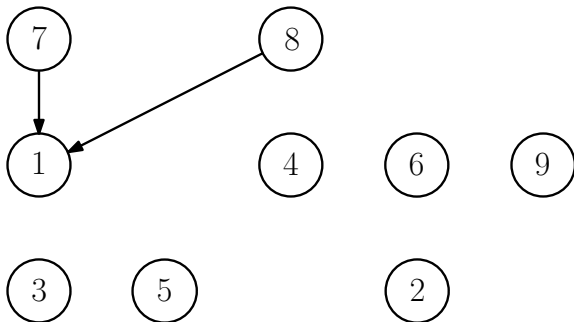
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

The algorithm is only allowed to erase part of its progress N times. If that happens, the algorithm shuffles A_C and B_C and starts over avoiding the choice leading to its first failed attempt.



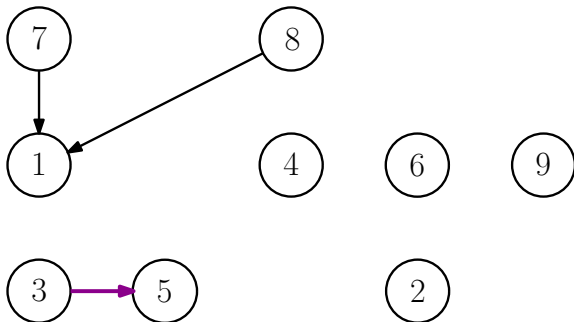
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



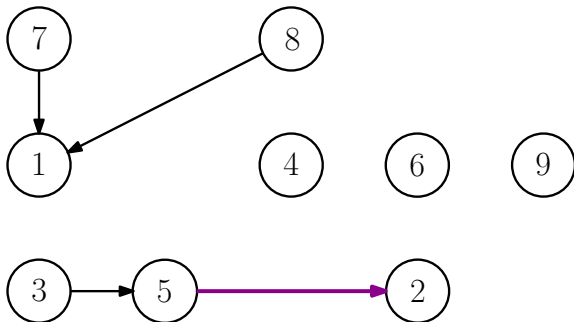
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



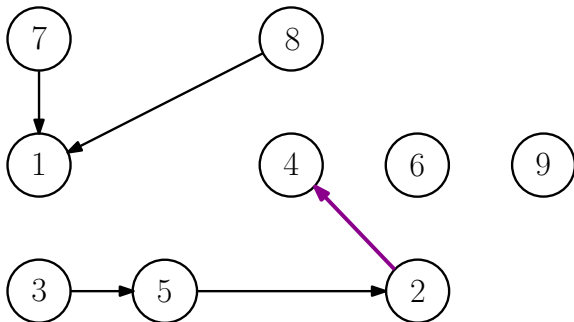
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



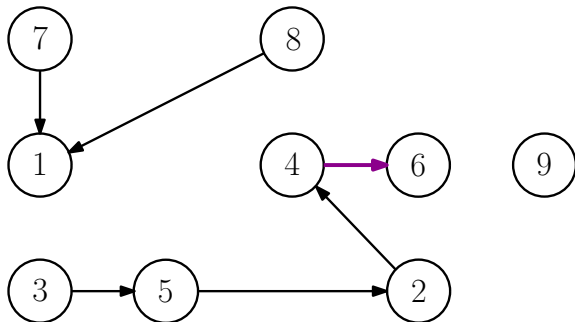
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

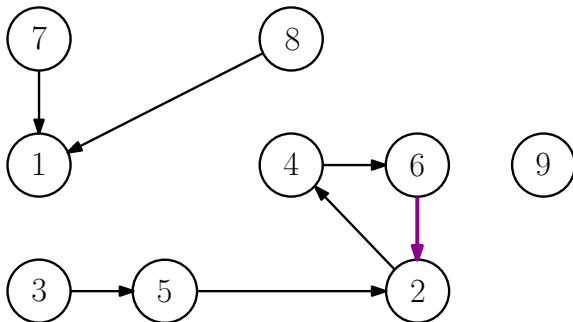
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

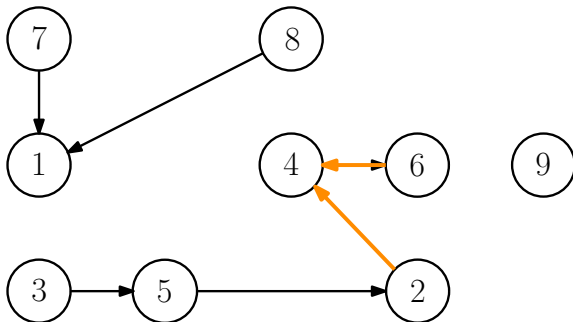
Unsafe edge (incompatible with (5, 2))! Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$
$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Problem resolved!



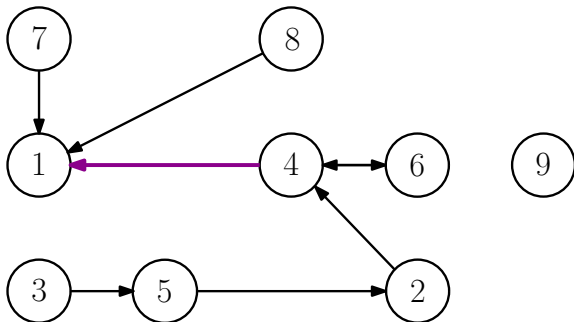
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



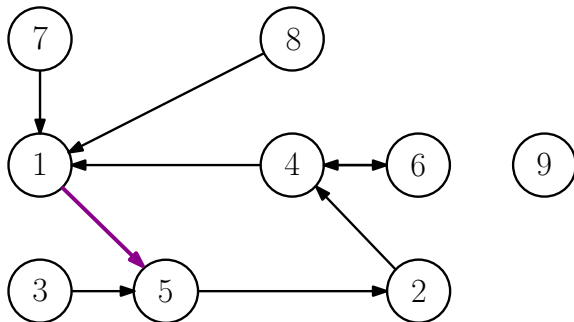
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



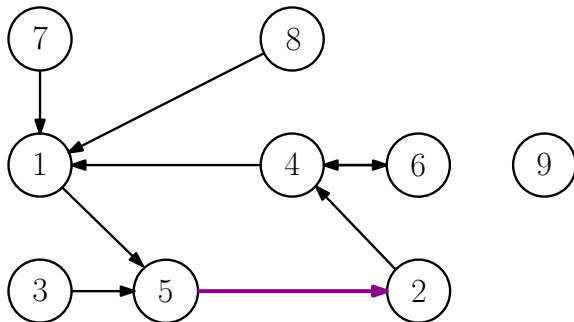
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

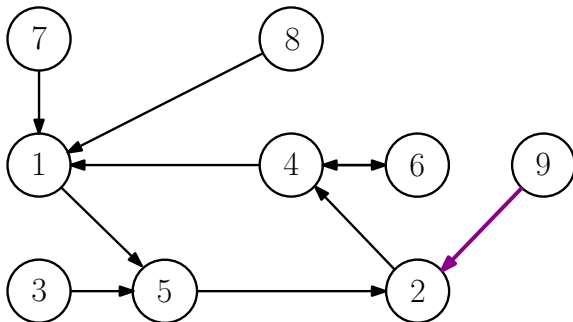
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Unsafe edge added
(incompatible with (5, 2))!
Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

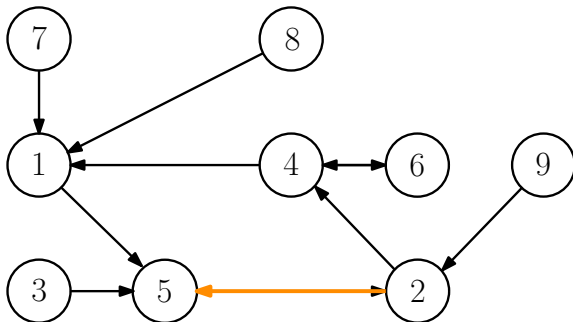
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Failed to remove (9, 2).
So, moved to (5, 2).
Failed to remove (5, 2).
So, flips (5, 2).
But now (2, 5) is
incompatible with (1, 5).



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

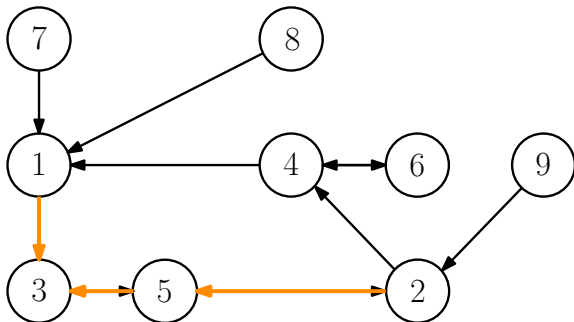
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

(9, 2) is still incompatible with another (5, 2) in the construction. So, the algorithm jumps back to the stage right before adding that (5, 2).



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

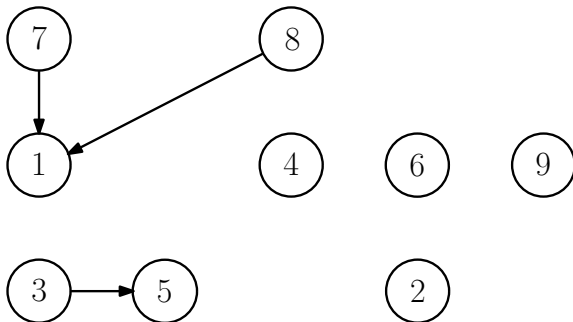
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

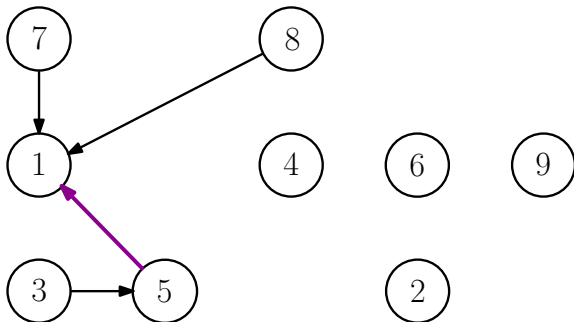
Shuffles A_C and B_C and starts over.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$
$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

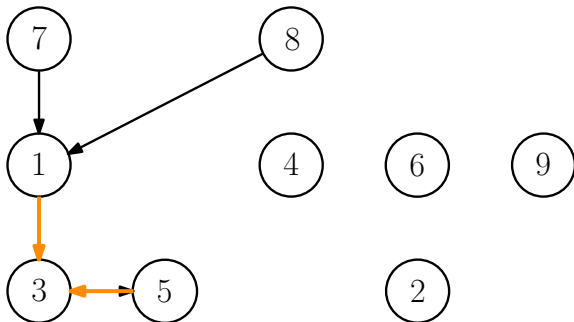
Unsafe edge (incompatible with (7, 1))! Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$
$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

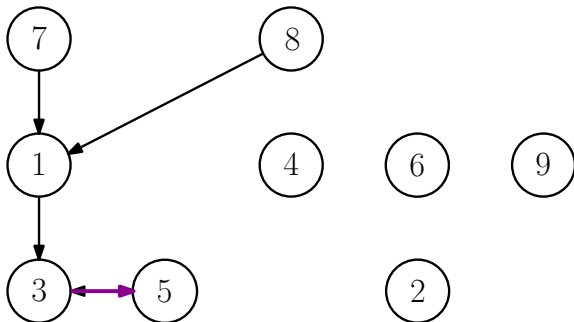
Problem resolved!



$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$


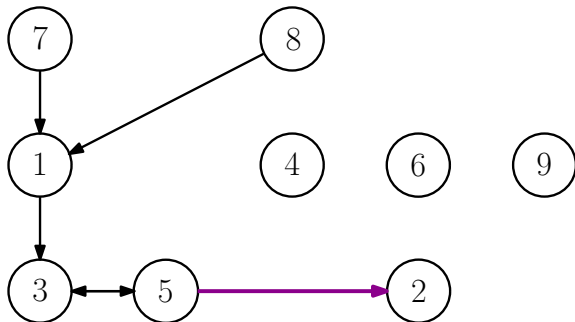
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



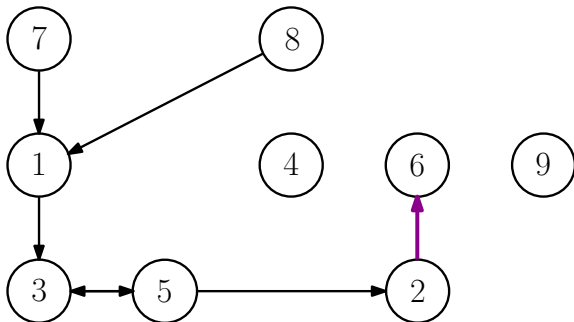
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



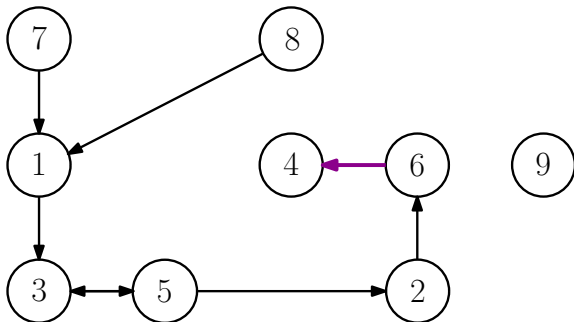
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



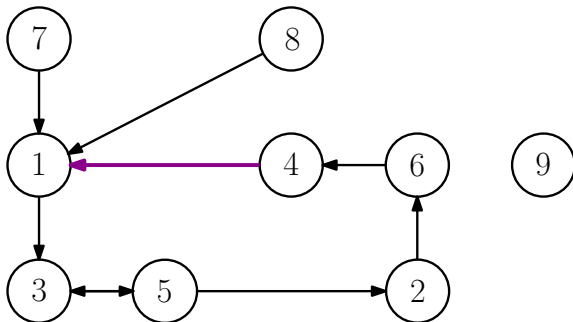
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

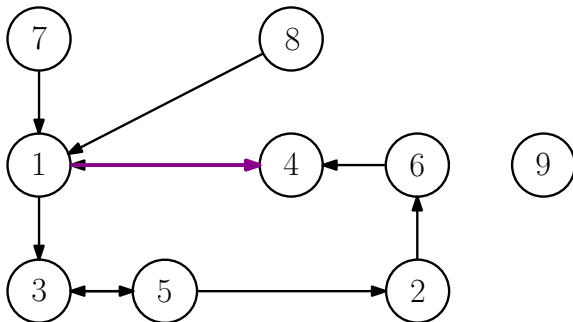
$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$
$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Unsafe edge (incompatible with (6, 4))! Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

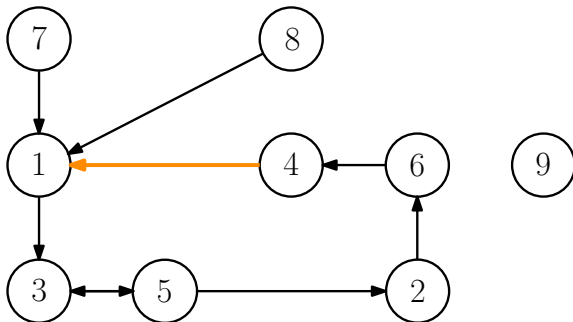
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

$$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\},$$

$$ComCh_C := \{(7, 8), (8, 7)\}, \quad NoComCh_C := \{(7, 9), (9, 7)\}.$$

Failed to remove (1, 4). So,
flips (1, 4) and now the
problem is resolved!



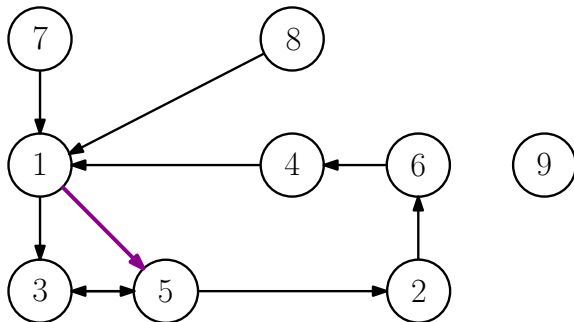
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



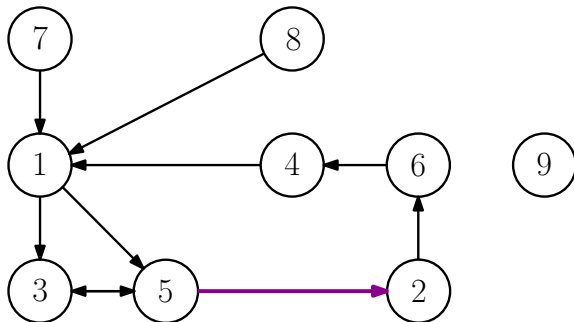
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



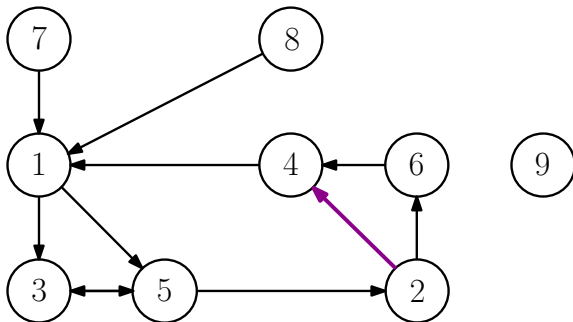
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$C := \{1, 2, 3, 4, 5, 6\}$,

$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\}$,

$B_C := \{(7, 1), (8, 1), (7, 4), (8, 4), (9, 6), (9, 2)\}$,

$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

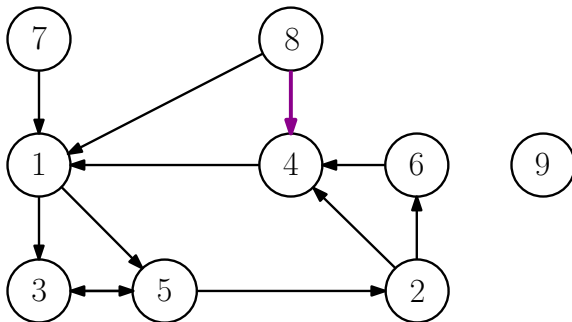
$$C := \{1, 2, 3, 4, 5, 6\},$$

$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$

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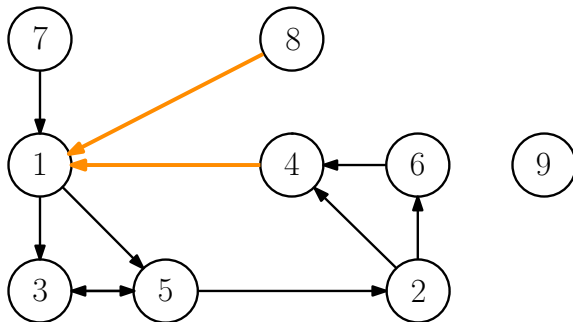
Unsafe edge (incompatible with (2, 4))! Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
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Problem resolved!



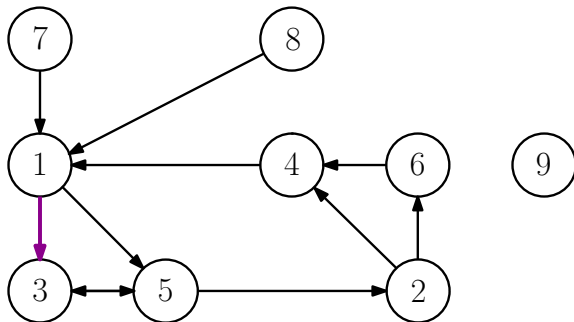
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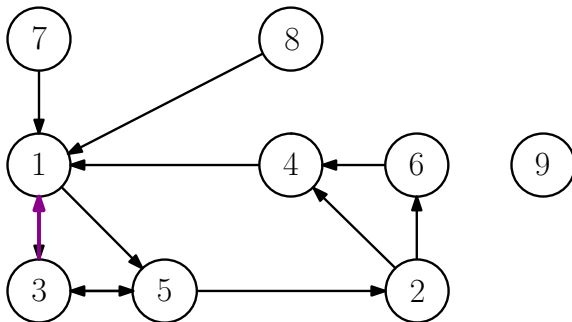
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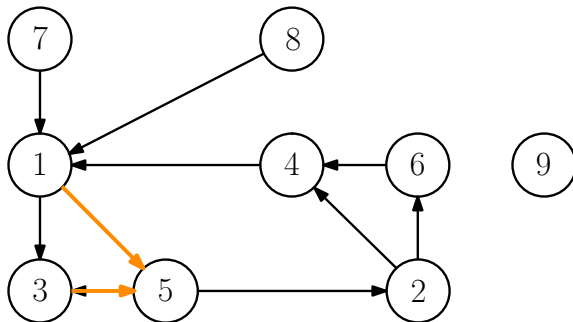
Unsafe edge (incompatible with (7, 1))! Correction process starts.



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$$C := \{1, 2, 3, 4, 5, 6\},$$
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Problem resolved!



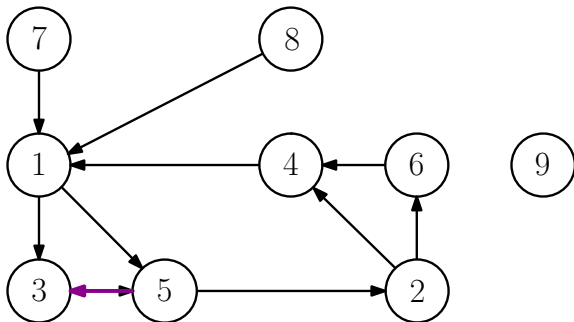
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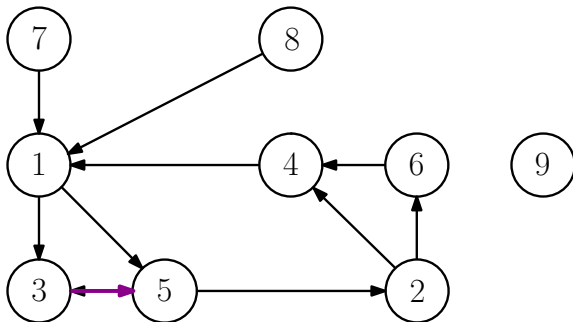
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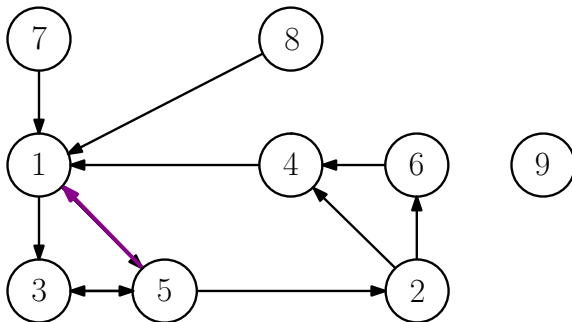
$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



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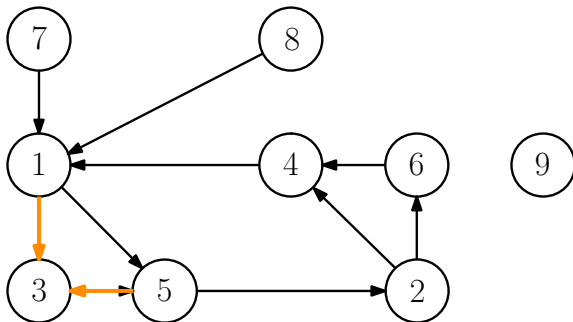
Unsafe edge (incompatible with (7, 1))! Correction process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
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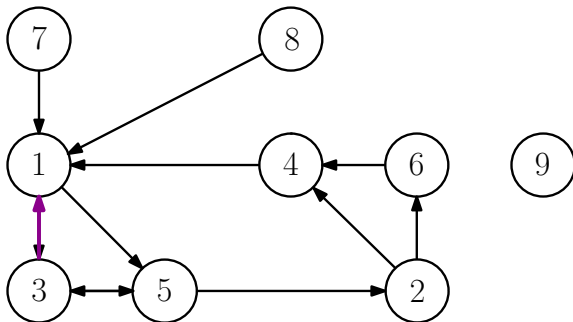
Problem resolved!



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
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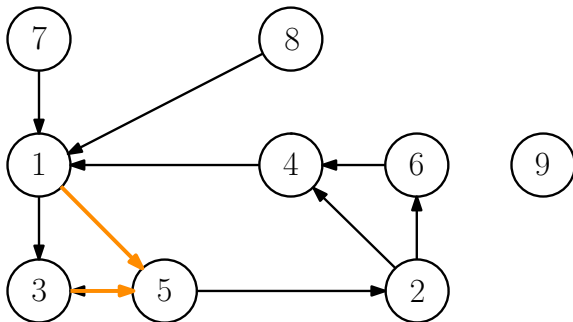
Unsafe edge (incompatible with (7, 1))! Correction process starts.



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$$C := \{1, 2, 3, 4, 5, 6\},$$
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Unsafe edge (incompatible with (7, 1))! Correction process starts.



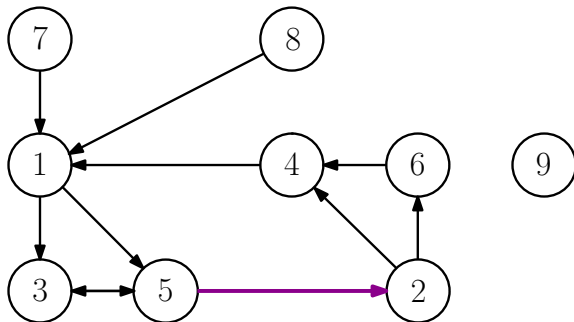
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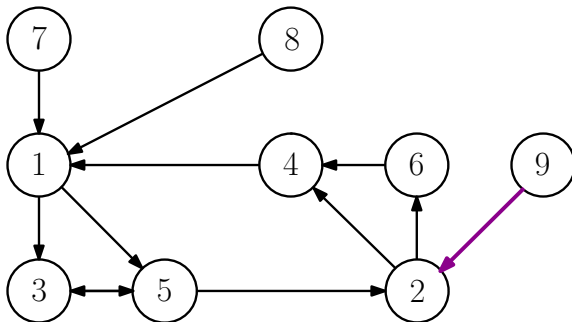
$ComCh_C := \{(7, 8), (8, 7)\}$, $NoComCh_C := \{(7, 9), (9, 7)\}$.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
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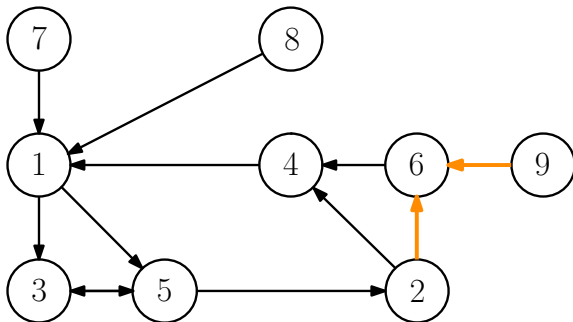
Unsafe edge (incompatible
with (5, 2))! Correction
process starts.



AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

$$C := \{1, 2, 3, 4, 5, 6\},$$
$$A_C := \{(1, 3), (3, 5), (1, 5), (2, 4), (2, 6), (4, 6), (1, 4), (2, 5)\},$$
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Problem resolved!



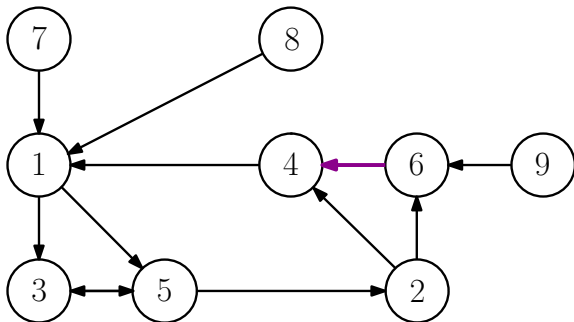
AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

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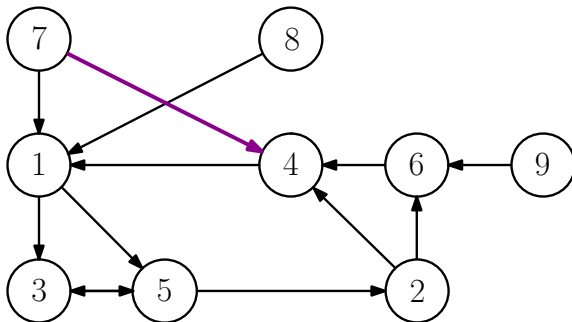
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Unsafe edge (incompatible
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AN SCCR ALGORITHM: CONSTRUCT AND CORRECT

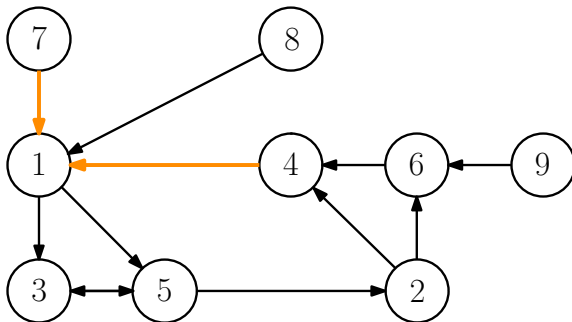
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Problem resolved!
The algorithm now outputs
this construction.



PROPOSITION

Suppose a subset $C \subseteq [n]$ and sets $A_C \subseteq C^2$, $B_C \subseteq ([n] \setminus C) \times C$, $ComCh_C \subseteq ([n] \setminus C)^2$ and $NoComCh_C \subseteq ([n] \setminus C)^2$ are given. For any $N \in \mathbb{N}$, if the SCCR algorithm outputs a set $E_C \subseteq [n]^2$, then E_C satisfies the desired properties.

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CONJECTURE

For all $n \in \mathbb{N}$, there exists $N \in \mathbb{N}$ that satisfies the following: Let (\mathcal{P}_0, π_0) be the partially ordered partition associated with G^* . Suppose $C \in \mathcal{P}_0$ and

$$A_C := \left\{ (a, b) \in E_{(\mathcal{P}_0, \pi_0)}^{(1)} \mid a, b \in C \right\},$$

$$B_C := \left\{ (a, b) \in E_{(\mathcal{P}_0, \pi_0)}^{(1)} \mid a \notin C, b \in C, C_{a, \mathcal{P}_0} \leq \pi_0 C_{b, \mathcal{P}_0} \right\}.$$

Moreover, let $ComCh_C$ and $NoComCh_C$ be the sets of pairs that must have and must not have common children in C per conditions of the previously stated theorem. Then given $C, A_C, B_C, ComCh_C, NoComCh_C$ and N , with high probability the SCCR algorithm outputs a set $E_C \subseteq [n]^2$ satisfying the desired properties.

PROPOSITION

Suppose a subset $C \subseteq [n]$ and sets $A_C \subseteq C^2$, $B_C \subseteq ([n] \setminus C) \times C$, $ComCh_C \subseteq ([n] \setminus C)^2$ and $NoComCh_C \subseteq ([n] \setminus C)^2$ are given. For any $N \in \mathbb{N}$, if the SCCR algorithm outputs a set $E_C \subseteq [n]^2$, then E_C satisfies the desired properties.

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[Back to simulations](#)